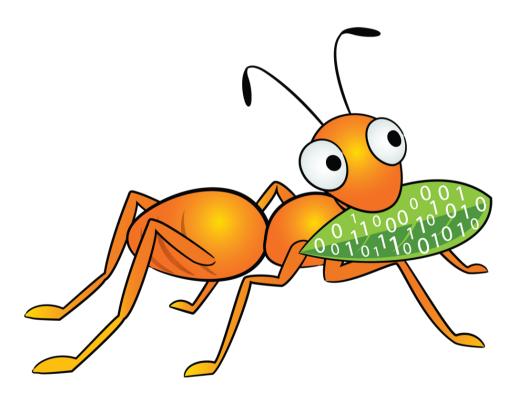
High-availability data storage and access on Gluster



Niels de Vos Red Hat Storage Engineer GlusterFS co-maintainer ndevos@redhat.com

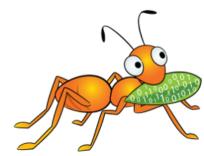


November 12, 2016 Eindhoven

Agenda

- What is Gluster?
 - Access Protocols/Methods
 - Basic High-Availability
- High-Availability projects
- Failure Scenarios
 - backend, frontend and applications
- Deployment Examples

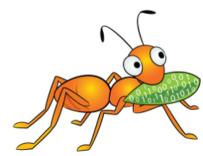




What is Gluster ?

Gluster is a distributed scale out filesystem that allows rapid provisioning of additional storage based on your storage consumption needs. It incorporates automatic failover as a primary feature. All of this is accomplished without a centralized metadata server.

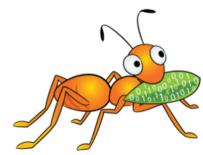




Gluster in Keywords

- Scalable, general-purpose storage platform
 - POSIX-y Distributed File System
 - Object storage (swift)
 - Flexible storage (libgfapi)
- No Metadata Server
- Heterogeneous Commodity Hardware
- Flexible and Agile Scaling
 - Capacity Petabytes and beyond
 - Performance Thousands of Clients

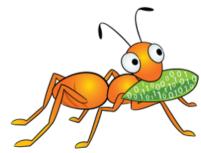




Gluster Use-Cases

- Content Delivery Networks
 - Media Streaming
 - Download Servers
- Archival
 - Backup services
 - Long term media archives
- Virtual Machine images
- High Performance / Distributed Computing
 - Rendering Farms

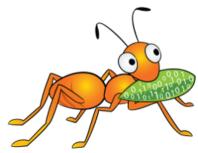




Data Access Overview

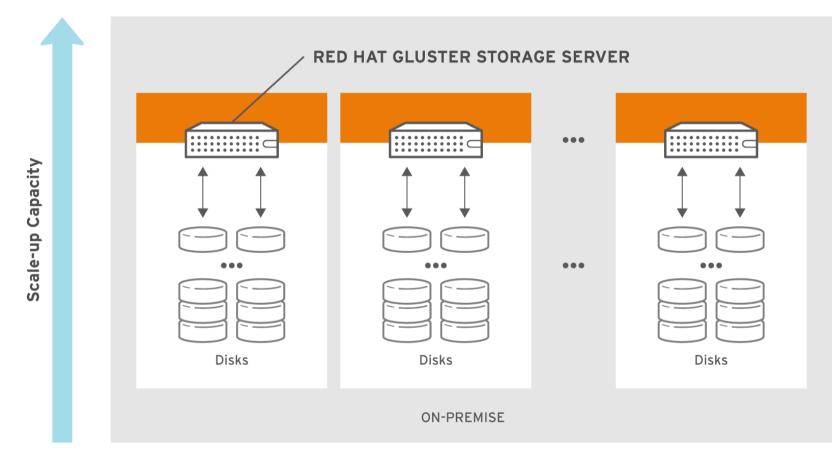
- GlusterFS Native Client
 - Filesystem in Userspace (FUSE)
- NFS
 - Built-in Service, NFS-Ganesha with libgfapi
- SMB/CIFS
 - Samba server required (libgfapi based module)
- Gluster For OpenStack Swift (Glusterswift)
- libgfapi flexible abstracted storage
 - Integrated with QEMU, Bareos and others





Scale-out and Scale-up

Scale-out performance, capacity and availability

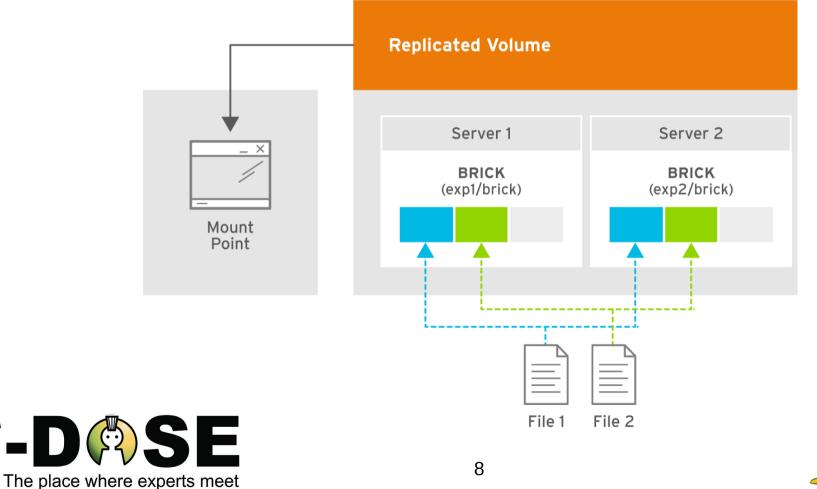


#145075_GLUSTER_1.0_334434_0415



Replicated Volume

- Copies files to multiple bricks
- Similar to file-level RAID 1





High-Availability projects

- CTDB from the Samba project
 - Used for Samba and Gluster/NFS
- Pacemaker from Cluster Labs
 - Used for NFS-Ganesha
 - Storhaug for common Samba + NFS-Ganesha
- keepalived
 - For high-availability + loadbalancing (webservers etc)





Failure Scenarios: backend

- Network failures within the storage communication
- Gluster clients can detect brick failures
- Gluster servers can detect storage-server failures
- Quorum is used for preventing split-brain scenarios
- Ping-timeout (client-side) for keeping connections active





Failure Scenarios: frontend

- NFS-Ganesha or Samba fail-over
- Standard HA-projects for IP-address relocation
 - Gluster integrated pacemaker with 10 sec. interval
- 'tickle' TCP-connections
- Potentially needs to notify clients for lock-reclaim
 - Gluster servers release locks after connection loss

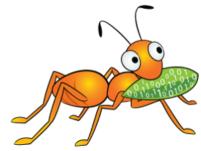




Failure Scenarios: applications

- Transparent for most applications
- Unless clients use file- or byte-range locks
 - When should a lock of a failing client be released?
- NFSv4 has a RENEW operation
- NFS-clients need to RECLAIM locks after failover





Deployment Examples: NFS failover

- NFS-Ganesha with NFSv4 clients
- Assume a NFS-Ganesha failure
- pacemaker detects the failure
 - By default a 10 second interval
- IP-address relocation
- NFS-Ganesha invokes GRACE on all servers
 - special state for NFS-server, by default 90 seconds
 - new lock requests (and I/O) get blocked/stalled
- NFS-clients need to reclaim their obtained locks **T-DOSE** The place where experts meet

Deployment Examples: Hypervisor and VMs

- oVirt management with QEMU/gfapi access to Gluster
- Assume a storage-server failure
- QEMU/gfapi detects the failure after ping-timeout
 - By default the ping-timeout is 42 seconds
- Filesystem in the VM may detect I/O delay sooner
 - By default SCSI emulation times-out after 30 seconds
- Problem: filesystems in VMs may become read-only



Deployment Examples: failover planning

- Assume 3 frontend servers (each has 33% of the load)
- One server fails, 33% of the load needs to be relocated
- One of the two remaining servers needs to handle 66%
 - Servers need double the resources for failover
- A problem when each server has only one virtual-IP
- Can be reduced when each server has two virtual-IPs
 - Servers need 1.5x the resources for failover









####